

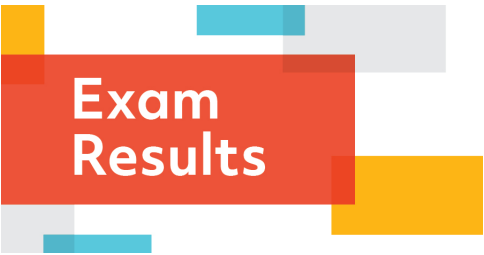
Dynamic Programming IV



Class 28

April 26, 2023

Probability Distribution for
Daily Demand
Assignment 10
Team Project 2



Exam Results

Probabilistic Dynamic Programming I



Henry Brewster decides to market the Fancy Assortment through three local outlets: Shaw's, Hannaford, and the Middlebury Co-Op.

Each day he produces 6 cases (24 boxes each) at a cost to him of \$100 a case. Each case sold at an outlet yields \$200.

Any unsold assortments are returned to the factory where he can sell them at \$50 a case as stale products the next day.

Probability Distribution for Daily Demand
(demand in cases)

		1	2	3
Store 1	Shaw's	.6	0	.4
Store 2	Hannaford's	.5	.1	.4
Store 3	Middlebury Co-Op	.4	.3	.3

Problem: How should he allocate the 6 cases to the three outlets to maximize his expected revenue?

Observations

- ▶ Don't give more than 3 cases to any store
- ▶ Distribute all 6

Expected Revenue Earned from Allocating x_n cases to store n
(in hundreds of dollars)

x_n	Store 1 Shaw's	Store 2 Hannaford's	Store 3 Co-Op
0	\$0	\$0	\$0
1	\$2	\$2	\$2
2	\$3.10	\$3.25	\$3.40
3	\$4.20	\$4.35	\$4.35

Allocate 2 to Shaw's (Store 1)
 $.6$ (Sell 1, Return 1) + $.4$ (Sell 2)
 $.6(2 + 1/2) + .4(4) = 1.5 + 1.6 = 3.1$

Allocate 3 to Shaw's
 $.6$ (Sell 1, return 2) + $.4$ (sell 3)
 $.6(2 + 1) + .4(6) = 1.8 + 2.4 = 4.2$

Allocate 2 to Hannaford's (Store 2)

$$\begin{aligned} &.5 (\text{Sell 1, return 1}) + .5(\text{sell 2}) \\ &.5(2 + 1/2) + .5(4) = 3.25 \end{aligned}$$

Allocate 3 to Hannaford's

$$\begin{aligned} &.5(\text{sell 1, return 2}) + .1(\text{sell 2, return 1}) + .4(\text{sell 3}) \\ &.5(2 + 1) + .1(4 + 1/2) + .4(6) \\ &1.5 + .45 + 2.4 = 4.35 \end{aligned}$$

Allocate 2 to Co-Op

$$\begin{aligned} &.4(\text{sell 1, return 1}) + .6(\text{sell 2}) \\ &.4(2 + 1/2) + .6(4) = 1 + 2.4 = 3.4 \end{aligned}$$

Allocate 3 to Co-Op

$$\begin{aligned} &.4(\text{sell 1, return 2}) + .3(\text{sell 2, return 1}) + .3(\text{sell 3}) \\ &.4(2 + 1) + .3(4 + ?) + .3(6) = 1.2 + 1.35 + 1.8 = 4.35 \end{aligned}$$

Expected Revenue Earned from Allocating x_n cases to store n
(in hundreds of dollars)

x_n	Store 1 Shaw's	Store 2 Hannaford's	Store 3 Co-Op
0	\$0	\$0	\$0
1	\$2	\$2	\$2
2	\$3.10	\$3.25	\$3.40
3	\$4.20	\$4.35	\$4.35

Let $r_i(s)$ = expected revenue of giving s cases to store i .
Then the recursive relationship is

$$f_3(s) = r_3(s)$$
$$f_n(s) = r_n(s) + f_{n+1}^*(s - s)$$
$$f_n^*(s) = \max_{k \leq s} [r_n(k) + f_{n+1}^*(s - k)]$$

$$f_3(s) = r_3(s)$$

$n = 3$	Co-Op					
x_3	0	1	2	3	f_3^*	x_3^*
s						
0	0	—	—	—	0	0
1	0	2	—	—	2	1
2	0	2	3.4	—	3.4	2
3	0	2	3.4	4.35	4.35	3

$$f_2^*(s) = \max_{k \leq s} [r_2(k) + f_3^*(s - k)]$$

$n = 2$	Hannaford					
$s \backslash x_2$	0	1	2	3	f_2^*	x_2^*
0	$0 + 0 = 0$	—	—	—	0	0
1	$0 + 2 = 2$	$2 + 0 = 2$	—	—	2	0 or 1
2	$0 + 3.4 = 3.4$	$2 + 2 = 4$	$3.25 + 0$	—	4	1
3	$0 + 4.35 = 4.35$	$2 + 3.4 = 5.4$	$3.25 + 2 = 5.25$	$4.35 + 0 = 4.35$	5.40	1
4	—	$2 + 4.35 = 6.35$	$3.25 + 3.4 = 6.65$	$4.35 + 2 = 6.35$	6.65	2
5	—	—	$3.25 + 4.35 = 7.60$	$4.35 + 3.4 = 7.75$	7.75	3
6	—	—	—	$4.35 + 4.35 = 8.70$	8.70	3

$$f_1^*(s) = \max_{k \leq s} [r_1(k) + f_2^*(s - k)]$$

$n = 1$	Shaw's					
x_1	0	1	2	3	f_1^*	x_1^*
s						
6	0 + 8.70 = 8.70	2 + 7.75 = 9.75	3.10 +6.65 = 9.75	4.2 + 5.40 = 9.60	9.75	1 or 2

Solutions

Store	Number of Cases	Number of Cases
Shaw's	1	2
Hannaford	3	2
Co-Op	2	2

**Expected Value of
Revenue: \$975**

Expected Value of Revenue: \$975



Best Possible Outcome

Sell all 6 cases → \$1200



Worst Possible Outcome

Sell 1 at each store

3 sales, 3 returns

Revenue = $3 \cdot \$200 + 3 \cdot \$50 = \$750$

**Probabilistic
Dynamic
Programming II**
Choose The Ideal Mate

Choosing The Ideal Mate



Choosing The Ideal Mate

N candidates each with values 0 - 1000, uniformly distributed.

Uniform Distribution: For all a and b with $0 \leq a < b \leq 1000$, the probability that a value lies in the interval $[a, b]$ is proportional to the length of the interval. In this case, the probability would be

$$\frac{b-a}{1000}.$$

If we choose at random a number in an interval $[a, b]$, its expected value is the average of a and b .

Candidates are chosen at random and presented to you one at a time.

You either Accept or Reject and move on to the next candidate.

Our Question: How should you proceed if you wish to maximize the expected value of the candidate accepted?

[Different Question: What is the best policy if you wish to maximize the probability of picking the best candidate among the N ?]

After all our online chats,
it's great to finally
meet you in person.

Same here.



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Stages : candidates

x_n : decision variable" accept or reject.

s : value of current candidate

Stage N : N th candidate has value uniformly distributed on $[0,1000]$ so the Expected Value is 500.

Stage $N - 1$: Consider $N - 1$ st Candidate:

Accept if $s > 500$; that is s is in the interval $[500,1000]$

Expected Value is $\frac{500+1000}{2} = 750$.

Reject is $s \leq 500$ in which case the Expected Value is 500

Thus **EV** = $\frac{1}{2}(750) + \frac{1}{2}(500) = \frac{1}{2}(750 + 500) = 625$

Stage $N - 2$: Consider $N - 2$ nd Candidate:

Accept if $s > 625$; that is, s is in the interval $[625, 1000]$

Expected Value is $\frac{625+1000}{2} = 812.50$

Reject is $s \leq 625$ in which case the Expected Value is 625

Thus **EV** = $(.625)(625) + (1 - .625)(812.50) = 695.3$

Stage $N - 3$: Consider $N - 3$ rd Candidate:

Accept if $s > 695.3$; that is, s is in the interval $[695.3, 1000]$

Expected Value is $\frac{695.3+1000}{2} = 847.5$

Reject is $s \leq 695.3$ in which case the Expected Value is 695.3

Thus **EV** = $(.6953)(695.3) + (1 - .6953)(847.50) = 741.5$

Stage $N - 3$: Consider $N - 3$ rd Candidate:

Accept if $s > 695.3$; that is, s is in the interval $[695.3, 1000]$

Expected Value is $\frac{695.3+1000}{2} = 847.5$

Reject is $s \leq 695.3$ in which case the Expected Value is 695.3

Thus **EV** = $(.6953)(695.3) + (1 - .6953)(847.50) = 741.5$

General Case

Accept if $s > x$; that is, s is in the interval $[x, 1000]$

Expected Value is $\frac{x+1000}{2}$

Reject is $s \leq x$ in which case the Expected Value is x

Thus **EV** = $\frac{x}{1000}x + (1 - \frac{x}{1000})(\frac{x+1000}{2})$

EV = $\frac{x^2}{1000} + (\frac{1000-x}{1000})(\frac{x+1000}{2})$

$$\mathbf{EV} = \frac{x^2}{1000} + \left(\frac{1000-x}{1000}\right)\left(\frac{x+1000}{2}\right)$$

$$\mathbf{EV} = \frac{2x^2}{2000} + \frac{1000^2-x^2}{2000} = \frac{x^2+1000^2}{2000}$$

$$\text{Improvement: } \frac{x^2+1000^2}{2000} - x =$$
$$\frac{x^2-2000x+1000^2}{2000} = \frac{(x-1000)^2}{2000}$$

Sequence: $x_1 = 500$; $x_n = \frac{x_n^2 + 1000^2}{2000}$

500

625

695

742

77

800

820

836

850

861

...

The sequence has a limit L which satisfies $L = \frac{L^2 + 1000^2}{2000}$
 $L^2 - 2000L + 1000^2 = 0 \rightarrow (L - 1000)^2 = 0 \rightarrow L = 1000$

What is the Incremental Improvement in Looking at One More Candidate?

$$\frac{(x-1000)^2}{2000}$$

$$x = 500 : \text{improvement is } \frac{(500-1000)^2}{2000} = \frac{250000}{2000} = 125$$

$$x = 850 : \text{improvement is } \frac{(850-1000)^2}{2000} = \frac{22500}{2000} = 11.25$$

When is the Incremental Improvement less than 5?

$$\frac{(x-1000)^2}{2000} < 5$$

$$(x - 1000)^2 > 5(2000) = 10000 = 10^4$$

$$|x - 1000| < 10^2$$

$$|x - 1000| < 100$$

AS SOON AS $x > 900$



**Next Time:
Use Dynamic Programming
To Solve
Fromage Cheese Company
Problem**