# Introduction to Game Theory III 

Class 31

May 3, 2023

# Finding Your Optimal Strategy in Zero-Sum Game 

- Eliminate All Dominated Strategies Determine "Best of the Worst"

|  | $T_{1}$ | $T_{2}$ | $T_{3}$ | $T_{4}$ | MIN |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :--- |
| $S_{1}$ | 5 | 6 | 20 | 3 | $\mathbf{3}$ |  |
| $S_{2}$ | 12 | 10 | 17 | 25 | $\mathbf{1 0}$ | $\Leftarrow$ |
| $S_{3}$ | 16 | 8 | 9 | 8 | $\mathbf{8}$ |  |
| $S_{4}$ | 13 | 9 | 6 | 5 | $\mathbf{5}$ |  |
| MAX | $\mathbf{1 6}$ | $\mathbf{1 0}$ | $\mathbf{2 0}$ | $\mathbf{2 5}$ |  |  |
|  |  | $\Uparrow$ |  |  |  |  |

maximin $=\underline{v}=$ lower value $\operatorname{minimax}=\bar{v}=$ upper value
If maximin $=$ minimax, then we have a Saddle Point.
If there is a Saddle Point, then it is stable. If both players know what the other will do, neither will change their strategy.

# NOT <br> EVERY GAME HAS A SADDLE POINT 

|  | $T_{1}$ | $T_{2}$ | $T_{3}$ | $T_{4}$ | MIN |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :--- |
| $S_{1}$ | 5 | 6 | 20 | 3 | $\mathbf{3}$ |  |
| $S_{2}$ | 12 | 10 | 17 | 25 | $\mathbf{1 0}$ | $\Leftarrow$ |
| $S_{3}$ | 16 | 8 | 9 | 8 | $\mathbf{8}$ |  |
| $S_{4}$ | 13 | 18 | 6 | 5 | $\mathbf{5}$ |  |
| MAX | $\mathbf{1 6}$ | $\mathbf{1 8}$ | $\mathbf{2 0}$ | $\mathbf{2 5}$ |  |  |

## A Simpler Example

|  | $T_{1}$ | $T_{2}$ | Worst |
| :---: | :---: | :---: | :---: |
| $S_{1}$ | 8 | -3 | -3 |
| $S_{2}$ | -5 | 4 | -5 |
| Worst | 8 | 4 |  |

Here $\underline{v}=-3$ and $\bar{v}=4$.
Rose can guarantee herself -3 by playing $S_{1}$.
Colin can limit her winnings to 4 .
Rose tentatively selects $S_{1}$ and Colin selects $T_{2}$.
But if Rose knows Colin will play $T_{2}$, she should switch to $S_{2}$.
Knowing this, Colin switches to $T_{1}$ but.. Then Rose switches to $S_{1}$ causing Colin to...

## An Historical Example



## World War II

Click for More Information:
Operation Lüttich and Operation Tractable, August 1944


Omar Bradley


Günther Von Kluge
October 30, 1882 - August 19, 1944
More Abut Kluge

Example: In August 1944 after the invasion of Normandy, the Allies broke out of their beachhead at Avranches, France and headed into the main part of the country. The German General von Kluge, commander of the ninth army, faced two options:

- $T_{1}$ : Stay and attack the advancing Allied armies.
- $T_{2}$ : Withdraw into the mainland and regroup.

Simultaneously, General Bradley, commander of the Allied ground forces faced a similar set of options regarding the German ninth army:

- $S_{1}$ Reinforce the gap created by troop movements at Avranches
- $S_{2}$ Send his forces east to cut-off a German retreat
- $S_{3}$ Do nothing and wait a day to see what the adversary did.

In real life, there were no pay-off values, however General Bradley's diary indicates the scenarios he preferred in order. There are six possible scenarios. Bradley ordered them from most to least preferable and using this ranking, we can construct the game matrix.

|  | Von Kluge's | Strategies | Row Min |  |
| :---: | :---: | :---: | :---: | :---: |
| Bradley's Strategies | Attack | Retreat |  |  |
| Reinforce Gap | 2 | 3 | 2 |  |
| Move East | 1 | 5 | 1 |  |
| Wait | 6 | 4 | 4 | $\Leftarrow \underline{v}$ |
| Column Max | 6 | 5 |  |  |

Notice that the maximin value of the rows is not equal to the minimax value of the columns. This is indicative of the fact that there is not a pair of strategies that form an equilibrium for this game.

|  | Von Kluge's | Strategies | Row Min |  |
| :---: | :---: | :---: | :---: | :---: |
| Bradley's Strategies | Attack | Retreat |  |  |
| Reinforce Gap | 2 | 3 | 2 |  |
| Move East | 1 | 5 | 1 |  |
| Wait | 6 | 4 | 4 | $\Leftarrow \underline{v}$ |
| Column Max | 6 | 5 |  |  |

Suppose that von Kluge plays his minimax strategy to retreat then Bradley would do better not to play his maximin strategy (wait) and instead move east, cutting of von Kluge's retreat, thus obtaining a payoff of $(5,-5)$. But von Kluge would realize this and deduce that he should attack, which would yield a payoff of (1, $-1)$. However, Bradley could deduce this as well and would know to play his maximin strategy (wait), which yields payoff $(6,-6)$. However, von Kluge would realize that this would occur in which case he would decide to retreat yielding a payoff of $(4,-4)$. The cycle then repeats.

## How Should Each Play The Game?

Choose a Strategy At Random! But With What Probability?


## Mixed Strategies

Let $p$ be the probability Rose plays $S_{1}$.
The $1-p$ is the probability she plays $S_{2}$.
Similarly, $q$ and $1-q$ would represent the probabilities that Colin selects $T_{1}$ and $T_{2}$, respectively.
Then the probabilities of the various outcomes are given by

|  | $q$ | $1-q$ |
| :---: | :---: | :---: |
| $p$ | $p q$ | $p(1-q)$ |
| $1-p$ | $(1-p) q$ | $(1-p)(1-q)$ |

with payoffs

|  | $q$ | $1-q$ |
| :---: | :---: | :---: |
| $p$ | 8 | -3 |
| $1-p$ | -5 | 4 |

Expected Value for Rose $=$
$\mathbf{E V}=8 p q-3 p(1-q)+(-5)(1-p) q+4(1-p)(1-q)$

## Expected Value for Rose =

$\mathbf{E V}=8 p q-3 p(1-q)+(-5)(1-p) q+4(1-p)(1-q)$

$$
\begin{gathered}
=20 p q-7 p-9 q+4 \\
=\left(4 p-\frac{9}{5}\right)\left(5 q-\frac{7}{4}\right)-\frac{63}{4}+4 \\
=\left(4 p-\frac{9}{5}\right)\left(5 q-\frac{7}{4}\right)+\frac{17}{20}
\end{gathered}
$$

$$
\text { Rose: } 4 p=\frac{9}{5} \Rightarrow p=\frac{9}{20}
$$

$$
\text { Colin: } 5 q-\frac{7}{4} \Rightarrow q=\frac{7}{20}
$$

|  | $T_{1}$ | $T_{2}$ | Worst |
| :---: | :---: | :---: | :---: |
| $S_{1}$ | 8 | -3 | -3 |
| $S_{2}$ | -5 | 4 | -5 |
| Worst | 8 | 4 |  |

Here $\underline{v}=-3$ and $\bar{v}=4$.

$$
E V=\left(4 p-\frac{9}{5}\right)\left(5 q-\frac{7}{4}\right)+\frac{17}{20}
$$

With $p=\frac{9}{20}$ and $q=\frac{7}{20}$, the Expected Value of the Game is $\frac{17}{20}$

## $\mathbf{E V}_{\text {Rose }}=20 p q-7 p-9 q+4$

$$
=(20 q-7) p-9 q+4
$$

So Colin should choose $q=\frac{7}{20}$
$\mathbf{E V}_{\text {Colin }}=-(20 p q-7 p-9 q+4)$

$$
\begin{gathered}
=-20 p q+7 p+9 q-4= \\
(9-20 p) q+7 p-4
\end{gathered}
$$

So Rose should choose $p=\frac{9}{20}$

## Another Way To Determine Equilibrium Strategies

Rose's View:
Mixed Strategy vs $T_{1}: 8 p+(1-p)(-5)=13 p-5$
Mixed Strategy vs $T_{2}:-3 p+4(1-p)=4-7 p$
These payoffs are equal when

$$
\begin{gathered}
13 p-5=4-7 p \\
20 p=9 \\
p=\frac{9}{20}
\end{gathered}
$$

Colin's Perspective:
Mixed Strategy vs $S_{1}: 8 q+(1-q)(-3)=11 q-3$
Mixed Strategy vs $S_{2}:-5 q+4(1-q)=4-9 q$
These payoffs are equal when

$$
11 q-3=4-9 q \Rightarrow 20 q=7 \Rightarrow q=\frac{7}{20}
$$

$$
\begin{aligned}
& \text { A Graphical Approach } \\
& \text { Look at Rose's Expected Payoffs: } \\
& \text { vs } T_{1}: 13 p-5 \\
& \text { vs } T_{2}: 4-7 p
\end{aligned}
$$

Graph these on $0 \leq p \leq 1$.


A plot of $13 p-5$ in red and $4-7 p$ in blue

Plot $\min (13 p-5,4-7 p)$ for $0 \leq p \leq 1$ -

Another Example

|  | $T_{1}$ | $T_{2}$ | $T_{3}$ | $T_{4}$ |
| :---: | :---: | :---: | :---: | :---: |
| $S_{1}$ | 1 | 0 | 3 | -3 |
| $S_{2}$ | -1 | 4 | -2 | 6 |


|  | $T_{1}$ | $T_{2}$ | $T_{3}$ | $T_{4}$ | Row Minima |
| :---: | :---: | :---: | :---: | :---: | :--- |
| $S_{1}$ | 1 | 0 | 3 | -3 | $\mathbf{- 3}$ |
| $S_{2}$ | -1 | 4 | -2 | 6 | $\mathbf{- 2}$ |
| Column | $\mathbf{1}$ | $\mathbf{4}$ | $\mathbf{3}$ | $\mathbf{6}$ |  |
| Maxima |  |  |  |  |  |


|  | $T_{1}$ | $T_{2}$ | $T_{3}$ | $T_{4}$ | Row Minima |
| :---: | :---: | :---: | :---: | :---: | :--- |
| $S_{1}$ | 1 | 0 | 3 | -3 | $-\mathbf{3}$ |
| $S_{2}$ | -1 | 4 | -2 | 6 | $-\mathbf{- 2} \Leftarrow \underline{v}$ |
| Column | $\mathbf{1}$ | 4 | 3 | 6 |  |
| Maxima | $\Uparrow$ |  |  |  |  |


|  | $T_{1}$ | $T_{2}$ | $T_{3}$ | $T_{4}$ | Row Minima |
| :---: | :---: | :---: | :---: | :---: | :--- |
| $S_{1}$ | 1 | 0 | 3 | -3 | $\mathbf{- 3}$ |
| $S_{2}$ | -1 | 4 | -2 | 6 | $\mathbf{- 2} \Leftarrow \underline{v}$ |
| Column | $\mathbf{1}$ | 4 | 3 | 6 |  |
| Maxima | $\Uparrow$ |  |  |  |  |
|  | $\bar{v}=1$ |  |  |  |  |

Value of this game is somewhere between -2 and 4.
Consider Expected Payoff to Rose if she uses $S_{1}$ with probability $p$ and $S_{2}$ with probability $1-p$.

$$
\begin{aligned}
& \text { vs } T_{1}: 1 p+(-1)(1-p)=2 p-1 \\
& \text { vs } T_{2}: 0 p+4(1-p)=4-4 p \\
& \text { vs } T_{3}: 3 p+(-2)(1-p)=5 p-2 \\
& \text { vs } T_{4}:-3 p+6(1-p)=6-9 p
\end{aligned}
$$

Expected Payoff to Rose with mixture $(p, 1-p)$ :

$$
\begin{aligned}
& \text { vs } T_{1}: 2 p-1 \\
& \text { vs } T_{2}: 4-4 p \\
& \text { vs } T_{3}: 5 p-2 \\
& \text { vs } T_{4}: 6-9 p
\end{aligned}
$$

Is there a single $p$ which guarantees same expected payoff against all 4 of Colin's strategies?
$T_{1}$ and $T_{2}: 2 p-1=4-4 p \Rightarrow 6 p=5 \Rightarrow p=\frac{5}{6}$
With $p=\frac{5}{6}$ :
Expected Value against $T_{1}=2\left(\frac{5}{6}\right)-1=2 / 3$
Expected Value against $T_{2}=4-4\left(\frac{5}{6}\right)=2 / 3$
But
Expected Value against $T_{3}=5\left(\frac{5}{6}\right)-2=13 / 6$
Expected Value against $T_{4}=6-9\left(\frac{5}{6}\right)=-3 / 2$

## Graphical Approach




Look at bottom edge


Look at bottom edge

Plot $\min (2 p-1,4-4 p, 5 p-2,6-9 p)$ for $0 \leq p \leq 1$


Rose can guarantee herself expected payoff of at least $\frac{3}{11}$ by choosing the strategy mixture $\left(\frac{7}{11}, \frac{4}{11}\right)$.
Colin asks: Can I keep her winnings down to $\frac{3}{11}$ ?


Colin wants to play a combination of $T_{1}$ and $T 4$.

|  | $q$ | $1-q$ |
| :---: | :---: | :---: |
|  | $T_{1}$ | $T_{4}$ |
| $S_{1}$ | 1 | -3 |
| $S_{2}$ | -1 | 6 |

Expected Payoffs

$$
\begin{aligned}
& \text { vs } S_{1}: 1 q-3(1-q)=4 q-3 \\
& \text { vs } S_{2}:-1 q+6(1-q)=6-7 q
\end{aligned}
$$

We can solve for $q$ by
Setting $4 q-3=\frac{3}{11}$ or
Setting $6-7 q=\frac{3}{11}$ or
Setting $4 q-3=6-7 q$
All these lead to $q=\frac{9}{11}$
Colin's Optimal Mixture is $\left(\frac{9}{11}, 0,0, \frac{2}{11}\right)$

We Can Use Graphical Approach Whenever One of the Players Has Exactly 2 Strategies

## Next Time: Connecting Game Theory With Linear Programming

