# Duality and The Fundamental Insight 

Class 12

March 10, 2023

## Handouts

- Cycling in Simplex Method Assignment 5


## Announcements



## Exam 1 Wednesday at 7 PM Warner 101

## Exam 1 Details

> No Time Limit
> Show All Work
> Double Check Your Answers Justify Claims Closed Book No Calculator

## What Are You "Responsible" For?

All Readings:

- Chapters 1-4 of Hillier and Lieberman
- Olinick LP Notes
- Cooper, Bhat and LeBlanc

Everything

- I've Said
- I Should Have Said But Didn't
- I Might Have Said
- The State of the World

Any questions?


## What Questions Remain?

Examine Final Tableaux of Fromage and Cheese Buyer Problems:

|  | Z | x | y | u | v | w |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | ---: |
| $Z$ | 1 | 0 | 0 | 0 | $5 / 12$ | $1 / 12$ | 1250 |
| $x$ | 0 | 1 | 0 | 0 | $1 / 6$ | $-1 / 6$ | 100 |
| $y$ | 0 | 0 | 1 | 0 | $-1 / 12$ | $5 / 24$ | 200 |
| $u$ | 0 | 0 | 0 | 1 | -4 | $5 / 2$ | 600 |

What is the meaning of the green numbers?

|  | $Z$ | $x$ | $y$ | $z$ | $u$ | $v$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z$ | 1 | 600 | 0 | 0 | 100 | 200 | -1250 |
| $y$ | 0 | 4 | 1 | 0 | $-1 / 6$ | $1 / 12$ | $5 / 12$ |
| $z$ | 0 | $-5 / 2$ | 0 | 1 | $1 / 6$ | $--5 / 24$ | $1 / 12$ |

## Duality And The

## Fundamental Insight

| Fromage Cheese Problem | Dual Problem |
| :---: | :---: |
| Maximize $4.5 x+4 y$ | Minimize $6000 C+2600 S+2000 B$ |
| subject to |  |
| $30 x+12 y \leq 6000$ (Cheddar) |  |
| $10 x+8 y \leq 2600$ (Swiss) |  |
| $4 x+8 y \leq 2000$ (Brie) | $30 C+10 S+4 B \geq 4.5$ |
| $x, y \geq 0$ | $12 C+8 S+8 B \geq 4$ |
| $x$ and $y$ are number of <br> packages of each assortment <br> to prepare | $C, S, B$ are the price per ounce <br> to offer for the cheeses |


|  | $x_{1}$ | $x_{2}$ |  |
| :--- | :---: | :---: | :--- |
| $y_{1}$ | 30 | 12 | $\leq 6000$ |
| $y_{2}$ | 10 | 8 | $\leq 2600$ |
| $y_{3}$ | 4 | 8 | $\leq 2000$ |
|  | IV | IV |  |
|  | 4.5 | 4 |  |

## Why Bother With Duality?

- Gain more insight into information the Simplex Method provides in the final tableau
- Run time (number of iterations) of Simplex Method depends more on the number of constraints than the number of decision variables. It may be quicker to solve the dual problem.

Example: Consider an LP with 10 variables and $5 \leq$ inequality constraints.
Convert to equations by introducing 5 slack variables.
We have a system of 5 equations in $10+5=15$ variables.
The number of potential bases is $\binom{15}{5}=3003$.
(a) Add 5 variables: 15 variables. 5 inequalities.

Number of potential bases is $\binom{20}{5}=15,504$.
(b) Add 5 new inequalities so we have 10 variables, $10 \leq$ constraints.
Number of potential bases is $\binom{20}{10}=184,756$.

## Final Tableaux

Tableau for the Optimal Basic Feasible Solution of Primal

|  | Z | x | y | u | v | w |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z$ | 1 | 0 | 0 | 0 | $5 / 12$ | $1 / 12$ | 1250 |
| $x$ | 0 | 1 | 0 | 0 | $1 / 6$ | $-1 / 6$ | 100 |
| $y$ | 0 | 0 | 1 | 0 | $-1 / 12$ | $5 / 24$ | 200 |
| $u$ | 0 | 0 | 0 | 1 | -4 | $5 / 2$ | 600 |

Tableau for the Optimal Basic Feasible Solution of the Dual

|  | Z | C | S | B | S1 | S2 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z$ | 1 | 600 | 0 | 0 | 100 | 200 | -1250 |
| $S$ | 0 | 4 | 1 | 0 | $-1 / 6$ | $1 / 12$ | $5 / 12$ |
| $B$ | 0 | $-5 / 2$ | 0 | 1 | $1 / 6$ | $-5 / 24$ | $1 / 12$ |

Fundamental Insight:

## Green Numbers Record Row Operations

Initial Tableau

|  | $Z$ | $x$ | $y$ | $u$ | $v$ | $w$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z$ | 1 | -4.5 | -4 | 0 | 0 | 0 | 0 |
| $u$ | 0 | 30 | 12 | 1 | 0 | 0 | 6000 |
| $v$ | 0 | 10 | 8 | 0 | 1 | 0 | 2600 |
| $w$ | 0 | 4 | 8 | 0 | 0 | 1 | 2000 |

Final Tableau

|  | $Z$ | $x$ | $y$ | $u$ | $v$ | $w$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z$ | 1 | 0 | 0 | $\mathbf{0}$ | $\mathbf{5} / \mathbf{1 2}$ | $\mathbf{1 / 1 2}$ | 1250 |
| $x$ | 0 | 1 | 0 | 0 | $1 / 6$ | $-1 / 6$ | 100 |
| $y$ | 0 | 0 | 1 | 0 | $-1 / 12$ | $5 / 24$ | 200 |
| $u$ | 0 | 0 | 0 | 1 | -4 | $5 / 2$ | 600 |

$$
\left(0, \frac{5}{12}, \frac{1}{12}\right) \cdot(6000,2600,2000)=\frac{0+13000+2000}{12}=\frac{15000}{12}=1250
$$

$$
\begin{gathered}
\left(\begin{array}{ccc}
0 & 1 / 6 & -1 / 6 \\
0 & -1 / 12 & 5 / 24 \\
1 & -4 & 5 / 2
\end{array}\right)\left(\begin{array}{l}
6000 \\
2600 \\
2000
\end{array}\right)=\left(\begin{array}{c}
100 \\
200 \\
600
\end{array}\right) \\
\left(\begin{array}{ccc}
0 & \frac{5}{12} & \frac{1}{12} \\
0 & 1 / 6 & -1 / 6 \\
0 & -1 / 12 & 5 / 24 \\
1 & -4 & 5 / 2
\end{array}\right)\left(\begin{array}{l}
C \\
S \\
B
\end{array}\right)=\left(\begin{array}{c}
\frac{5}{12} S+\frac{1}{12} B \\
\frac{S-B}{6} \\
\frac{1}{12} S+\frac{5}{24} B \\
C-4 S+\frac{5}{2} B
\end{array}\right)
\end{gathered}
$$

## A Peek Into Sensitivity Analysis

Suppose initial amounts of cheeses are
$C=6000, S=3000, B=2760$.
Then

$$
\left(\begin{array}{ccc}
0 & \frac{5}{12} & \frac{1}{12} \\
0 & 1 / 6 & -1 / 6 \\
0 & -1 / 12 & 5 / 24 \\
1 & -4 & 5 / 2
\end{array}\right)\left(\begin{array}{c}
6000 \\
3000 \\
2760
\end{array}\right)=\left(\begin{array}{c}
1480 \\
40 \\
325 \\
900
\end{array}\right)
$$

Optimal solution is $x=40, y=325, u=900$, Revenue $=1480$.

## Shadow Prices

$$
\begin{gathered}
(C, S, B)=\left(0, \frac{5}{12}, \frac{1}{12}\right) \text { are } \\
\text { Shadow Prices }
\end{gathered}
$$

They not only give us the optimal solution to the dual problem; they also tell us how much each additional ounce of cheese would boost the objective function in the primal problem.

| Primal Problem | Dual Problem |
| :---: | :---: |
| Maximize $Z=c \mathbf{x}$ | Minimize $W=\mathbf{y} b$ |
| subject to | subject to |
| $A \mathbf{x} \leq b$ | $\mathbf{y} A \geq c$ |
| and $\mathbf{x} \geq 0$ | and $\mathbf{y} \geq 0$. |

Note: If you like to write decision variables (unknowns) on the right, then you can write $A^{T} \mathbf{y}^{\top} \geq c^{T}$ instead of $\mathbf{y} A \geq c$

| Primal Problem | Dual Problem |
| :---: | :---: |
| Maximize $Z=c x$ | Minimize $W=y b$ |
| subject to | subject to |
| $A x \leq b$ | $y A \geq c$ |
| and $x \geq 0$. | and $y \geq 0$. |
| $\mathbf{x}:$ column vector | $\mathbf{y}:$ row vector |


| Alternative 1 for Dual | Alternative 2 for Dual |
| :---: | :---: |
| Minimize $W=\mathbf{b}^{T} \mathbf{w}$ | Maximize $W=-\mathbf{b}^{T} \mathbf{w}$ |
| subject to | subject to |
| $A^{T} \mathbf{w} \geq \mathbf{c}^{T}$ | $-A^{T} \mathbf{w} \leq-\mathbf{c}^{T}$ |
| and $\mathbf{w} \geq 0$ | and $\mathbf{w} \geq 0$. |
| $\mathbf{w}:$ column vector | $\mathbf{w}$ : column vector |

Theorem: The dual of the dual is the primal.

## Dual of the Dual is the Primal

| Primal | Alternative 2 for Dual |
| :---: | :---: |
| Maximize $Z=\mathbf{c x}$ | Maximize $W=-\mathbf{b}^{T} \mathbf{w}$ |
| subject to | subject to |
| $A \mathbf{x} \leq \mathbf{b}^{T}$ | $-A^{T} \mathbf{w} \leq-\mathbf{c}^{T}$ |
| and $\mathbf{x} \geq 0$ | and $\mathbf{w} \geq 0$. |
| $\mathbf{x}:$ column vector | $\mathbf{w}:$ column vector |

